

Jets, our window on partons at the LHC

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Partons — quarks and gluons — are key concepts of QCD.

- ▶ Lagrangian is in terms of quark and gluon fields
- ▶ Perturbative QCD *only* deals with partons
- ▶ Concept of parton powerful even beyond perturbation theory
 - hadron classifications
 - exotic states, e.g. colour glass condensate (high gluon densities)

Yet it is surprisingly hard to give unambiguous meaning to partons.

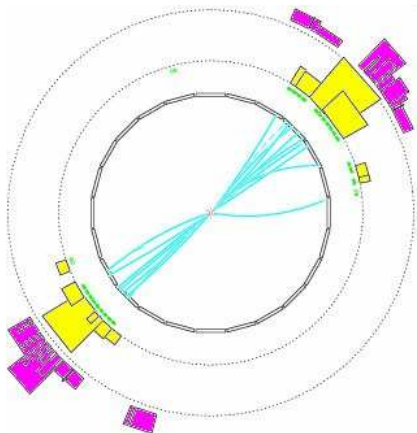
- ▶ Not an asymptotic state of the theory — because of confinement
- ▶ But also even in perturbation theory
 - because of collinear divergences (in massless approx.)

Despite this, there are two decent ways of “seeing” partons:

- ▶ Scatter some hard probe off them, e.g. a virtual photon → DIS
- ▶ See traces of them in the final state → jets

In each case ill-defined nature of a parton translates into ambiguity in the partonic interpretation of what you see

richness of the physics



Jets are what we see.

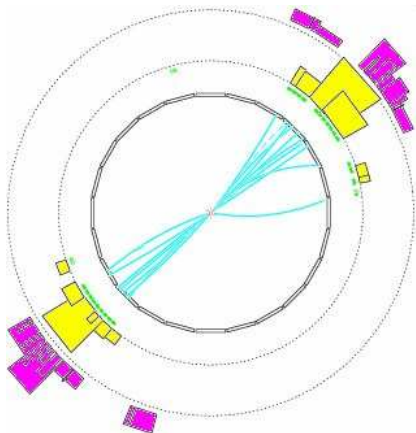
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$$E_{parton} = M_Z/2?$$

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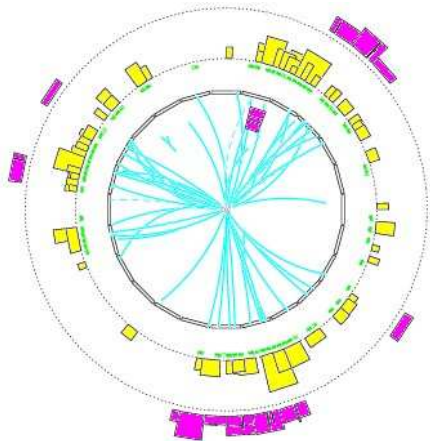
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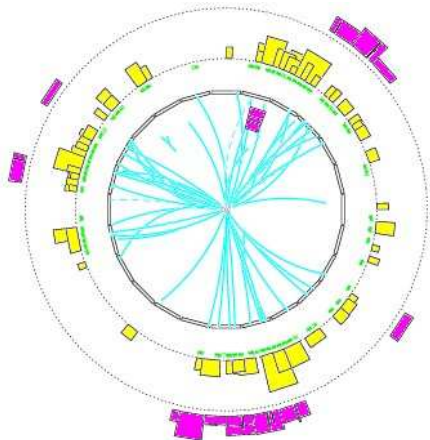
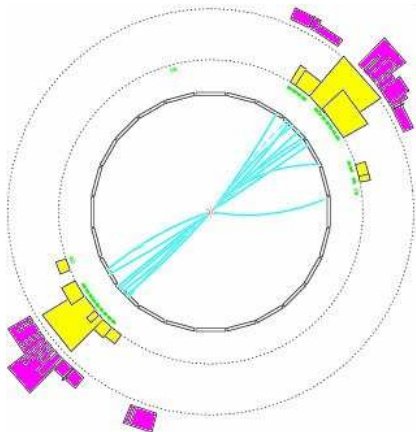
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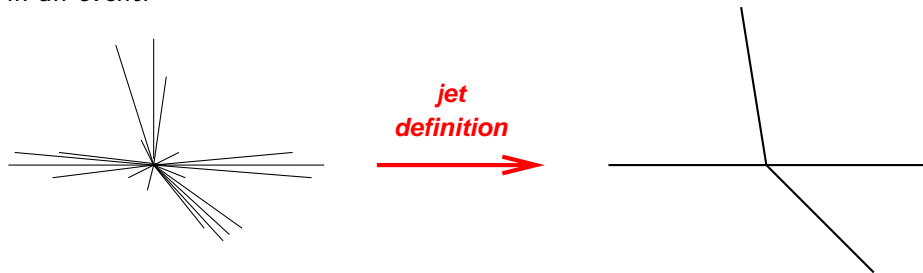
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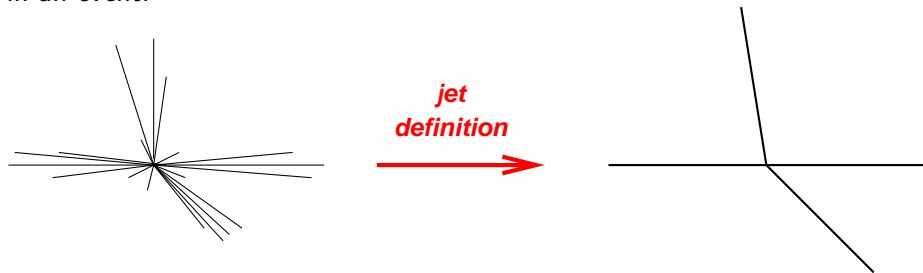


Jets are *as close as we can get to a physical single hard quark or gluon*: with good definitions their properties (multiplicity, energies, [flavour]) are

- ▶ finite at any order of perturbation theory
- ▶ insensitive to the parton \rightarrow hadron transition

NB: finiteness \longleftrightarrow set of jets depends on jet def.

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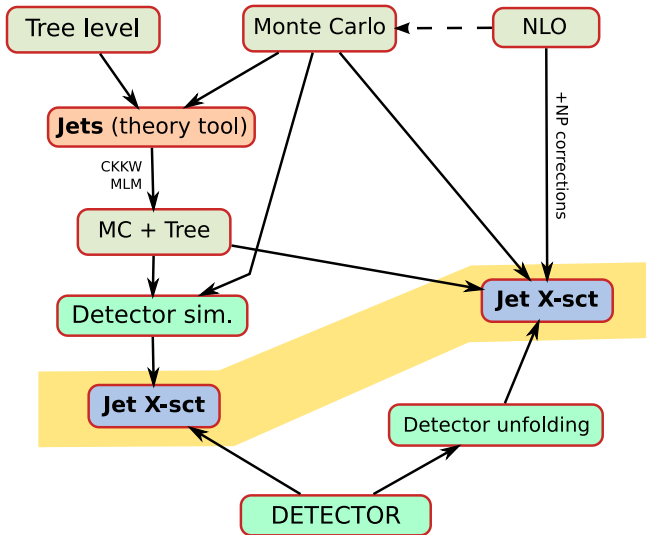
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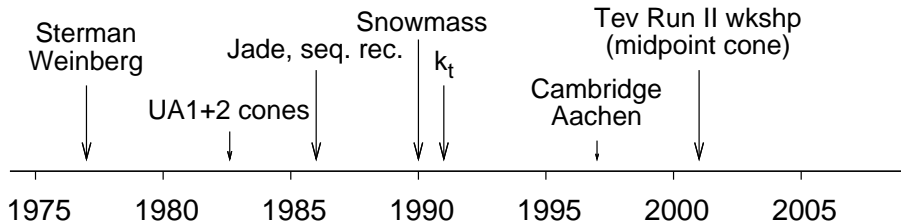
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Jet (definitions) provide central link between expt., "theory" and theory

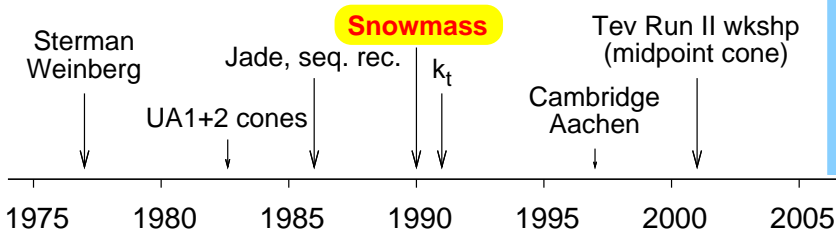
- ▶ Periodic key developments in jet definitions spurred by ever-increasing experimental/theoretical sophistication.
- ▶ Approach of LHC provides motivation for taking a new, fresh, systematic look at jets.
- ▶ This talk: **some of the discoveries along the way**



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NB: also ARCLUS, OJF, ...

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Speed, IR safety, Jet Areas
Non-pert. effects, Jet Flavour

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Experiment	N
LEP, HERA	50
Tevatron	100–400
LHC low-lumi	800
LHC high-lumi	4000
LHC PbPb	30000

- ▶ Range & complexity of signatures (jets, $t\bar{t}$, tj , Wj , Hj , $t\bar{t}j$, WWj , Wjj , SUSY, etc.)
- ▶ e.g. ~ 5 million $t\bar{t} \rightarrow 6$ jet events/year
- ▶ Theory investment
 ~ 100 people \times 10 years
 60 – 100 million \$

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Snowmass Accord (1990):

FERMILAB-Conf-90/249-E
[E-741/CDF]

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Several important properties that should be met by a jet definition are [3]:

1. Simple to implement in an experimental analysis;
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k_t , Jade, Cam/Aachen, ...

Bottom-up:

Cluster 'closest' particles repeatedly until few left → jets.

Works because of mapping:

closeness \Leftrightarrow *QCD divergence*

Loved by e^+e^- , ep and theorists

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UA1, JetClu, Midpoint, ...

Top-down:

Find coarse regions of energy flow (cones), and call them jets.

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Both had serious issues that got in way of practical use and/or physical validity

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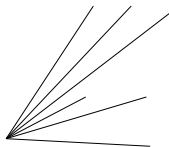
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- ▶ Find smallest of all $d_{ij} = \min(k_{ti}^2, k_{tj}^2) \Delta R_{ij}^2 / R^2$ and $d_{iB} = k_i^2$
- ▶ Recombine i, j (if iB : $i \rightarrow \text{jet}$)
- ▶ Repeat



NB: hadron collider variables

▶ $\Delta R_{ij}^2 = (\phi_i - \phi_j)^2 + (y_i - y_j)^2$

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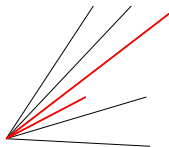
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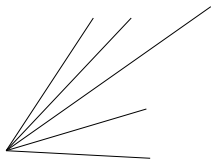
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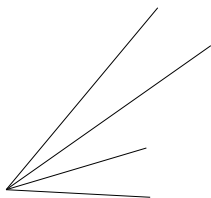
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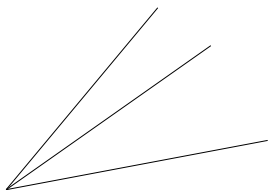
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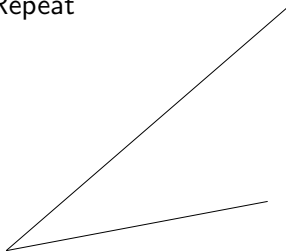
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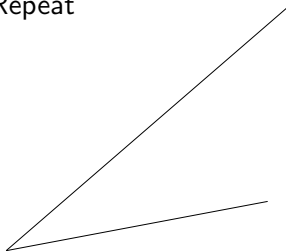
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are closely related to structure of divergences for QCD emissions

$$[dk_j] |M_{g \rightarrow g_i g_j}^2(k_j)| \sim \frac{\alpha_s C_A}{2\pi} \frac{dk_{tj}}{\min(k_{ti}, k_{tj})} \frac{d\Delta R_{ij}}{\Delta R_{ij}}, \quad (k_{tj} \ll k_{ti}, \Delta R_{ij} \ll 1)$$

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'Trivial' computational issue:

- ▶ for N particles: $N^2 d_{ij}$ searched through N times = N^3
- ▶ 4000 particles (or calo cells): **1 minute**
NB: often study $10^7 - 10^8$ events (20-200 CPU years)
- ▶ Heavy Ions: 30000 particles: **10 hours/event**

As far as possible physics choices should not be limited by computing.

Even if we're clever about repeating the full search each time, we still have $\mathcal{O}(N^2)$ d_{ij} 's to establish

Fast Hierarchical Clustering and Other Applications of Dynamic Closest Pairs

David Eppstein
UC Irvine

We develop data structures for dynamic closest pair problems with arbitrary distance functions, that do not necessarily come from any geometric structure on the objects. Based on a technique previously used by the author for Euclidean closest pairs, we show how to insert and delete objects from an n -object set, maintaining the closest pair, in $O(n \log^2 n)$ time per update and $O(n)$ space. With quadratic space, we can instead use a quadtree-like structure to achieve an optimal time bound, $O(n)$ per update. We apply these data structures to hierarchical clustering, greedy matching, and TSP heuristics, and discuss other potential applications in machine learning, Gröbner bases, and local improvement algorithms for partition and placement problems. Experiments show our new methods to be faster in practice than previously used heuristics.

Categories and Subject Descriptors: F.2.2 [Analysis of Algorithms]: Nonnumeric Algorithms

General Terms: Closest Pair, Agglomerative Clustering

Additional Key Words and Phrases: TSP, matching, conga line data structure, quadtree, nearest neighbor heuristic

1. INTRODUCTION

Hierarchical clustering has long been a mainstay of statistical analysis, and clustering based methods have attracted attention in other fields: computational biology (reconstruction of evolutionary trees; tree-based multiple sequence alignment), scientific simulation (n -body problems), theoretical computer science (network design and nearest neighbor searching) and of course the web (hierarchical indices such as Yahoo). Many clustering methods have been devised and used in these applications, but less effort has gone into algorithmic speedups of these methods.

In this paper we identify and demonstrate speedups for a key subroutine used in several clustering algorithms, that of maintaining closest pairs in a dynamic set of objects. We also describe several other applications or potential applications of the

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Eppstein '99
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Of these naive methods, brute force recomputation may be most commonly used, due to its low space requirements and ease of implementation. Three hierarchical clustering codes we examined, Zupan's [Zupan 1982], CLUSTAL W [Thompson et al. 1994], and PHYLIP [Felsenstein 1995] use brute force. (Indeed, they do not even save space by doing so, since they all store the distance matrix.) Pazzani's learning code [Pazzani 1997] also uses brute force (M. Pazzani, personal communication), as does *Mathematica*'s Gröbner basis code (D. Lichtblau, personal communication).

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- ▶ Consider smallest $d_{ij} = \min(k_{ti}^2, k_{tj}^2) R_{ij}^2$
- ▶ Suppose $k_{ti} < k_{tj}$
- ▶ Then: $R_{ij} \leq R_{i\ell}$ for any $\ell \neq j$. [If $\exists \ell$ s.t. $R_{i\ell} < R_{ij}$ then $d_{i\ell} < d_{ij}$]

In words: if i, j form smallest d_{ij} then j is geometrical nearest neighbour (GNN) of i .

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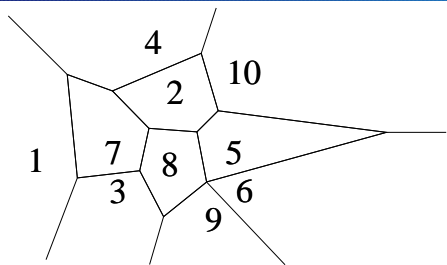
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Finding Geom Nearest Neighbours



Given a set of vertices on plane (1...10) a *Voronoi diagram* partitions plane into cells containing all points closest to each vertex

Dirichlet '1850, Voronoi '1908

A vertex's nearest other vertex is always in an adjacent cell.

E.g. GNN of point 7 will be found among 1,4,2,8,3 (it turns out to be 3)

Construction of Voronoi diagram for N points: $N \ln N$ time Fortune '88

Update of 1 point in Voronoi diagram: $\ln N$ time

Devillers '99 [+ related work by other authors]

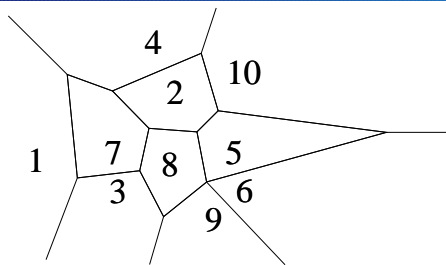
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Finding Geom Nearest Neighbours



Given a set of vertices on plane (1...10) a *Voronoi diagram* partitions plane into cells containing all points closest to each vertex

Dirichlet '1850, Voronoi '1908

A vertex's nearest other vertex is always in an adjacent cell.

E.g. GNN of point 7 will be found among 1,4,2,8,3 (it turns out to be 3)

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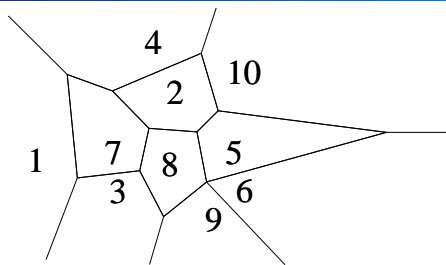
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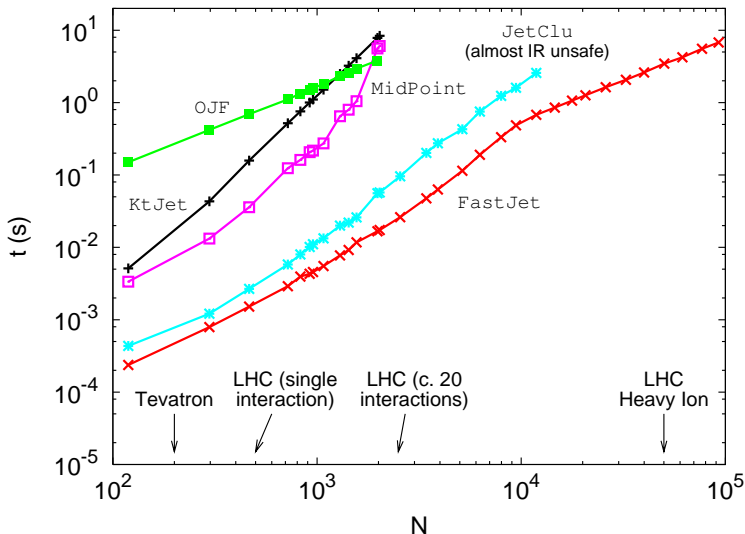
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NB: for $N < 10^4$, FastJet switches to a related geometrical N^2 alg.

Conclusion: speed issues for k_t resolved

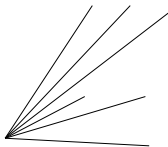
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▶ Find some/all stable cones

≡ cone pointing in same direction as the momentum of its contents

▶ Resolve cases of overlapping stable cones

By running a 'split-merge' procedure



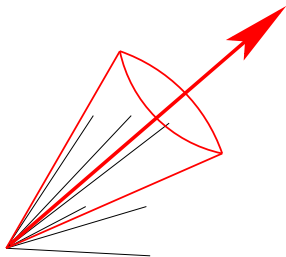
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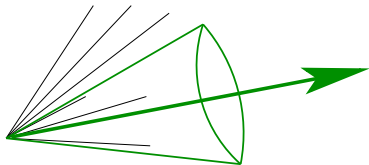
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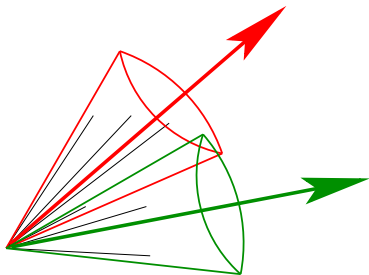
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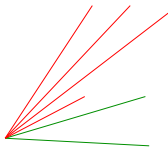
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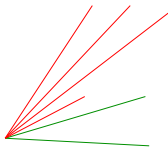
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Qu: How do you find the stable cones?

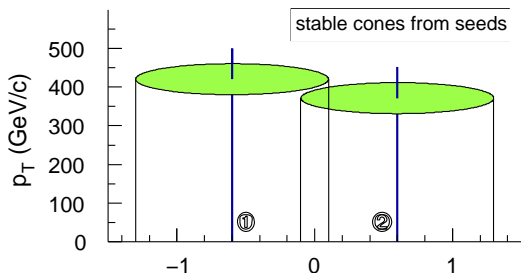
All experiments use iterative methods:

- ▶ use each particle as a starting direction for cone; use sum of contents as new starting direction; repeat.
- ▶ use additional 'midpoint' starting points between pairs of initial stable cones.

'Midpoint' algorithm



Use of seeds is *dangerous*



Extra soft particle adds new seed \rightarrow changes final jet configuration.

This is **IR unsafe**.

Divergences of real and virtual contributions do not cancel at $\mathcal{O}(\alpha_s^4)$

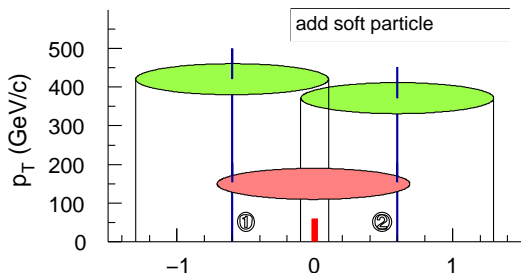
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Solution: add extra seeds at midpoints of all pairs, triplets, ... of stable cones.

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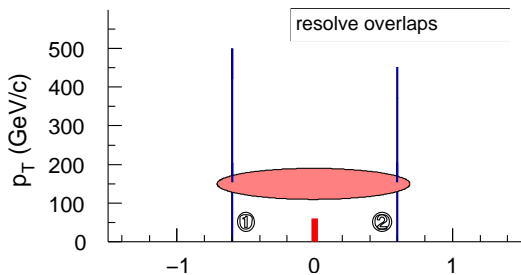
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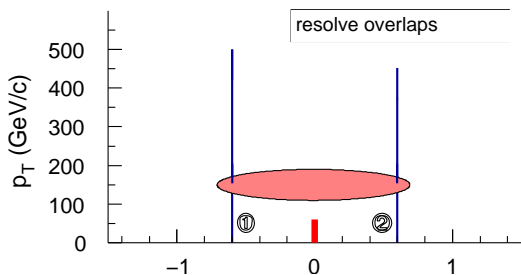
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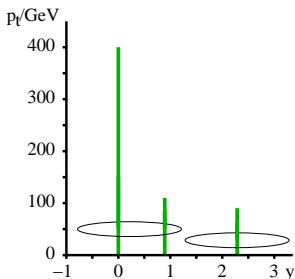
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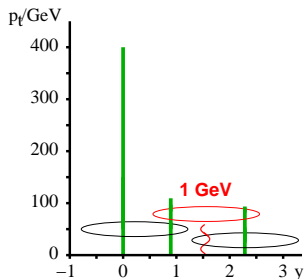
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Midpoint IR problem



Stable cones
with midpoint:

$\{1,2\}$ & $\{3\}$



$\{1,2\}$ & $\{2,3\}$ & $\{3\}$

Jets with
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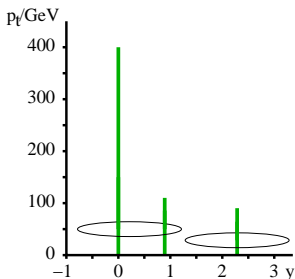
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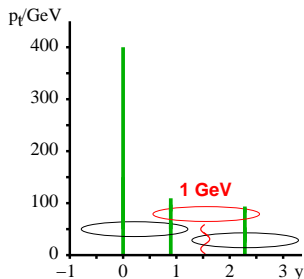
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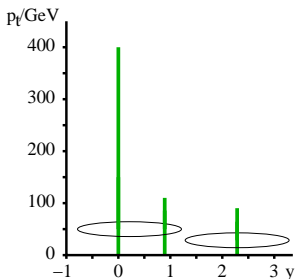
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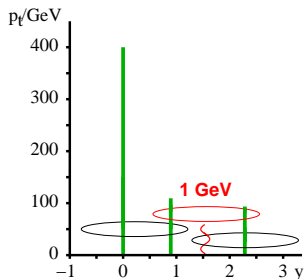
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Midpoint was supposed to solve *just this type of problem*. But worked only at lowest order.

IR/Collinear unsafety is a serious problem!

- ▶ Invalidates theorems that ensure finiteness of perturbative QCD
 - Cancellation of real & virtual divergences
- ▶ Destroys usefulness of (intuitive) partonic picture
 - you cannot think in terms of hard partons if adding a 1 GeV gluon changes 100 GeV jets
- ▶ 'Pragmatically:' limits accuracy to which it makes sense to calculate

Process	1st miss cones @	Last meaningful order
Inclusive jets	NNLO	NLO [NNLO being worked on]
W/Z + 1 jet	NNLO	NLO
3 jets	NLO	LO [NLO in nlojet++]
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A cone algorithm should find **all** stable cones

First advocated: Kidonakis, Oderda & Sterman '97

Guarantees IR safety of the set of stable cones

Only issue: you still need to find the stable cones in practice.

One known exact approach:

- ▶ Take each possible subset of particles and see if it forms a stable cone.
Tevatron Run II workshop, '00 (for fixed-order calcs.)
- ▶ There are 2^N subsets for N particles. Computing time $\sim N2^N$.
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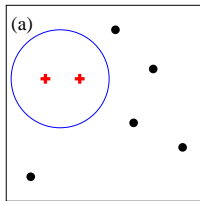
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Transform into a geometrical problem

Cones are just *circles* in the $y - \phi$ plane. To find all stable cones:

1. Find all distinct ways of enclosing a subset of particles in a $y - \phi$ circle
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Finding all distinct circular enclosures of a set of points is *geometry*:



Any enclosure can be moved until a pair of points lies on its edge.

Polynomial time recipe for finding all distinct enclosures:

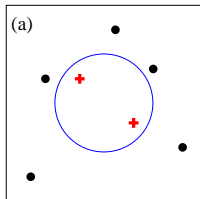
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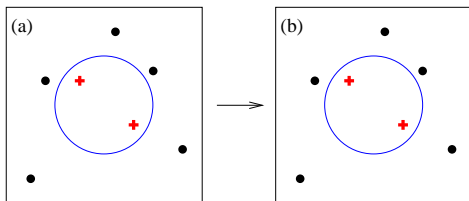
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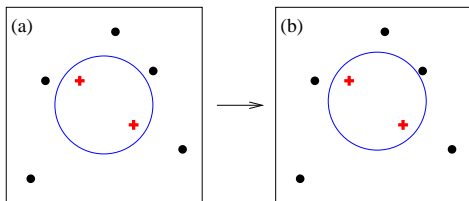
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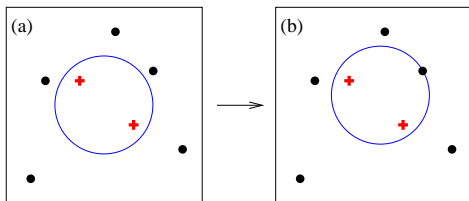
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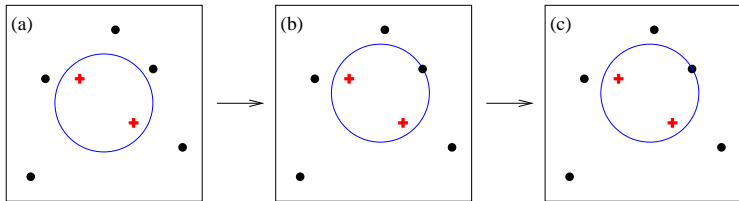
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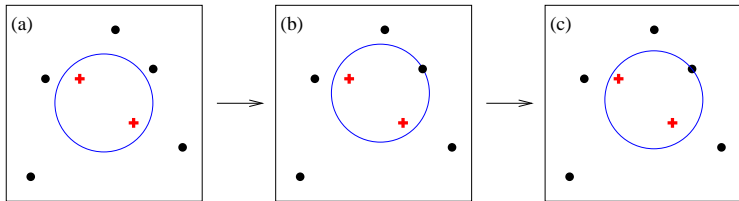
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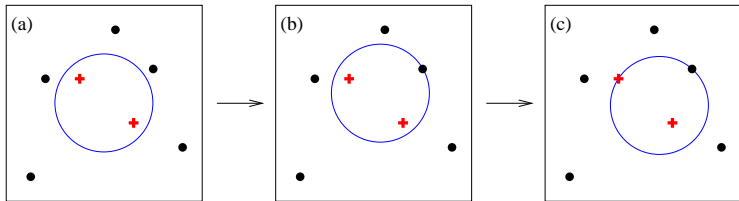
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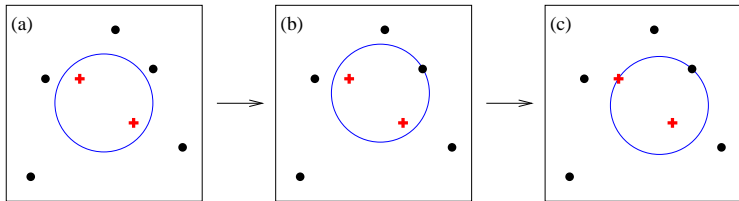
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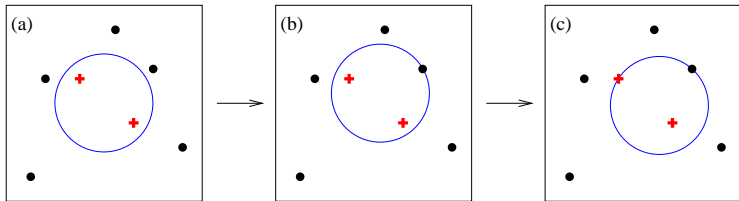
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Naive implementation of this idea would run in N^3 time.

N^2 pairs of points, pay N for each pair to check stability

N^3 is also time taken by midpoint codes (smaller coeff.)

With some thought, this reduces to $N^2 \ln N$ time.

Traversal order, stability check

checkxor

GPS & Soyez '07

- ▶ Much faster than midpoint with no seed threshold

IR unsafe

- ▶ Same speed as midpoint codes with seeds > 1 GeV

Collinear unsafe

- └ 2. Safe, practical jet-finding
- └ 2. Cone algorithms

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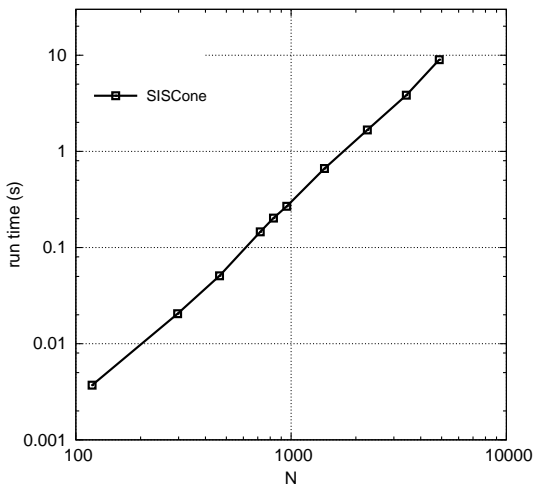
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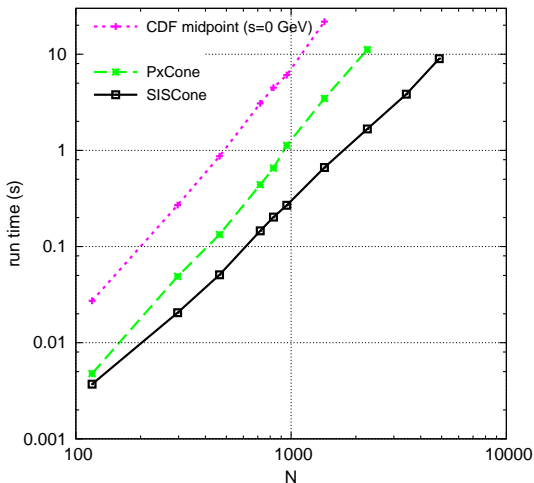
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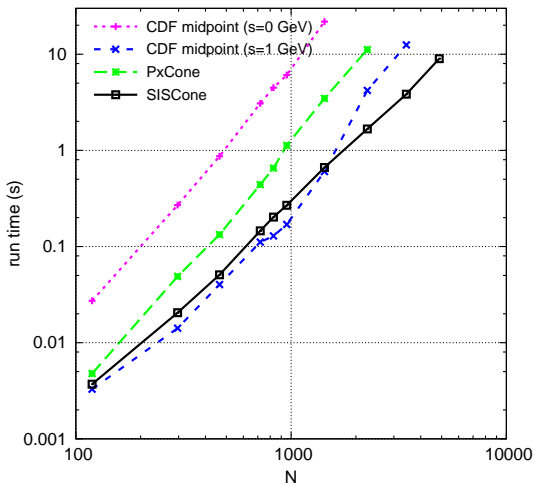
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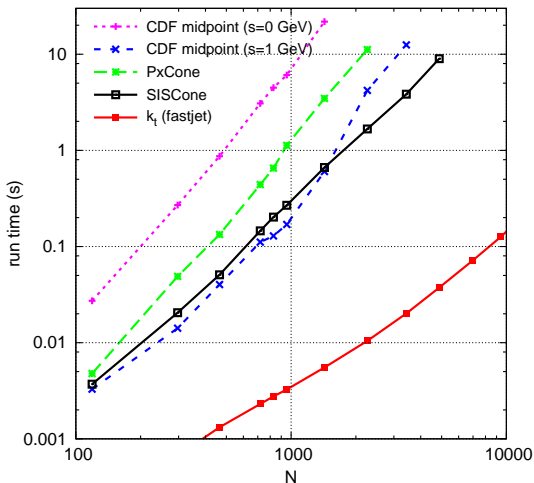
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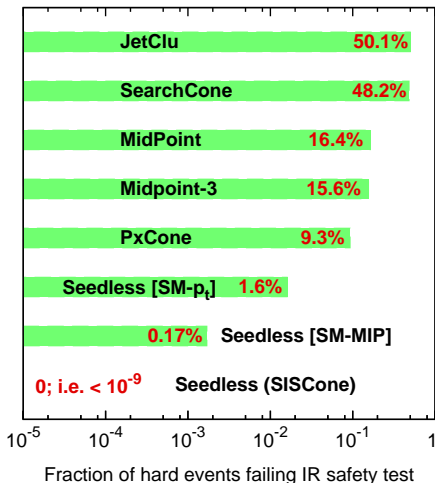
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Complementary set of IR/Collinear safe jet algs \longrightarrow flexibility in studying complex events.

Consider families of jet algs: e.g. sequential recombination with

$$d_{ij} = \min(k_{ti}^{2p}, k_{tj}^{2p}) \Delta R_{ij}^2 / R^2$$

	Alg. name	Comp. Geometry problem	time
$p = 1$	k_t CDOSTW '91-93; ES '93	Dynamic Nearest Neighbour CGAL (Devillers et al)	$N \ln N$ exp.
$p = 0$	Cambridge/Aachen Dok, Leder, Moretti, Webber '97 Wengler, Wobisch '98	Dynamic Closest Pair T. Chan '02	$N \ln N$
$p = -1$	anti- k_t (cone-like) Cacciari, GPS, Soyez, in prep.	Dynamic Nearest Neighbour CGAL (worst case)	$N^{3/2}$
cone	SISCone GPS Soyez '07 + Tevatron run II '00	All circular enclosures previously unconsidered	$N^2 \ln N$ exp.

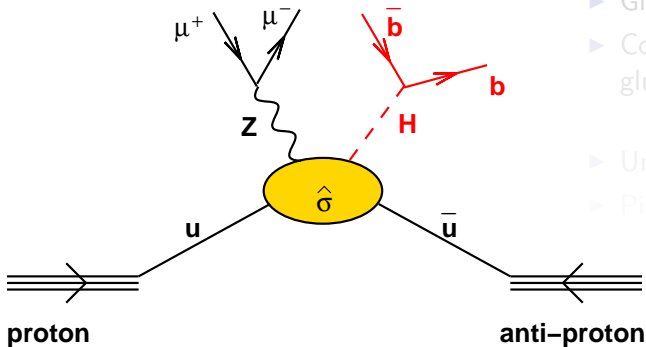
All accessible in FastJet
FastJet in software of all (4) LHC collaborations

Once you have a decent set of jet algs, *start asking questions about them.*

- ▶ They share a common parameter R (angular reach). How do results depend on R ?
- ▶ In what way do the various algorithms differ?
- ▶ How are they to be best used in the challenging LHC environment?

Try to answer questions with Monte Carlo? Gives little understanding of underlying principles.

➡ *Supplement with analytical approximations.*



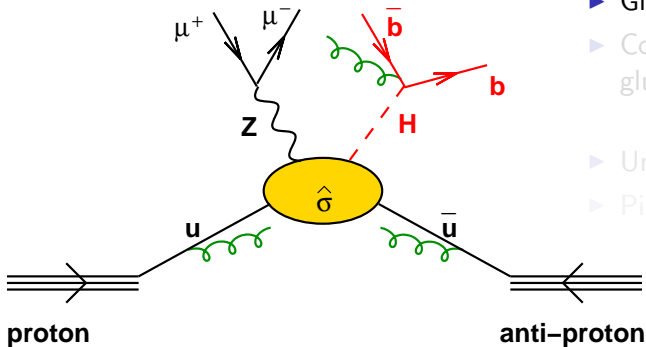
▶ Gluon emission, $\mathcal{O}(\alpha_s)$

▶ Conversion of quarks, gluons $\rightarrow \pi^\pm$, etc.

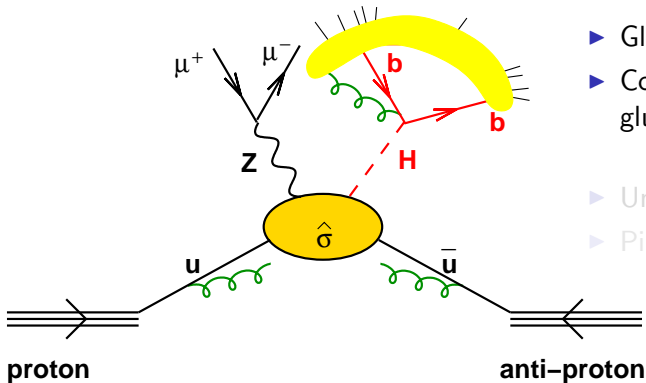
Hadronisation

▶ Underlying event

▶ Pileup



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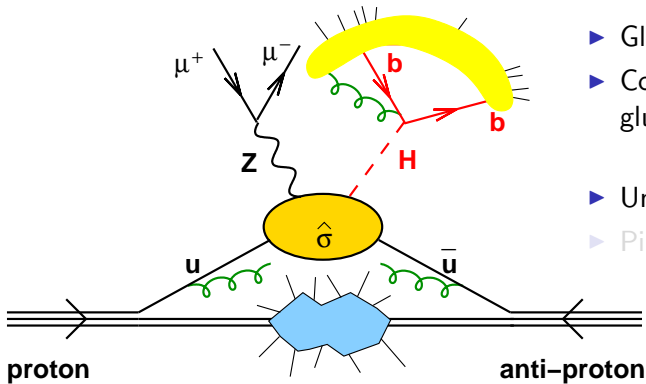
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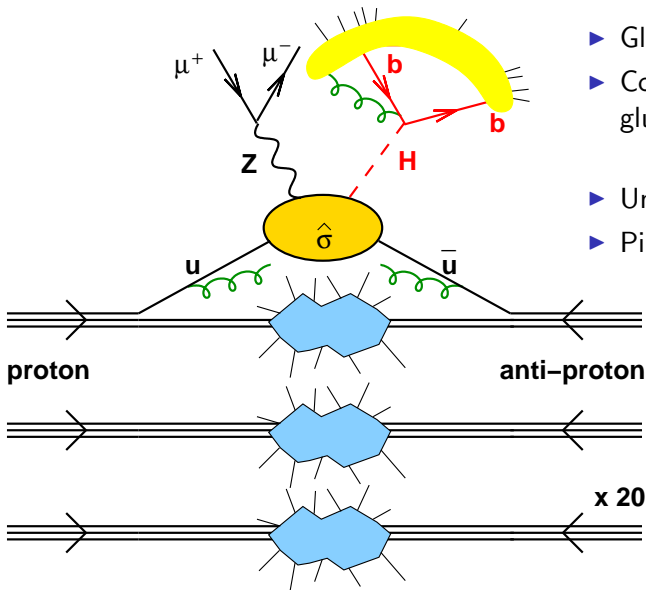
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Start with *quark* with transverse momentum p_t

$$\begin{aligned}
 \langle \delta p_t \rangle_{PT} &\simeq \frac{1}{\sigma_0} \int d\Phi |M^2| \alpha_s(k_{t,rel}) (p_{t,jet} - p_t) \\
 &\simeq \frac{\alpha_s C_F}{\pi} \int_R^{\mathcal{O}(1)} \frac{d\theta}{\theta} \int dz p_{gq}(z) \cdot ((1-z)p_t - p_t) \\
 &\simeq -1.01 \frac{\alpha_s C_F}{\pi} p_t \ln \frac{1}{R} + \mathcal{O}(\alpha_s p_t) \qquad C_F = 4/3
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Similarly for gluon:

$$\langle \delta p_t \rangle_{PT} \simeq -(0.94 C_A + 0.15 n_f T_R) \frac{\alpha_s}{\pi} p_t \ln \frac{1}{R} + \mathcal{O}(\alpha_s p_t) \qquad C_A = 3$$

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Dokshitzer & Webber; Korchemsky & Sterman

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$$\frac{2}{\pi} \delta\alpha_s(k_{t,rel}) = \Lambda \delta(k_{t,rel} - \Lambda)$$

$\Lambda = \int dk_{t,rel} \delta\alpha_s(k_{t,rel})$, should be
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Tested for ~ 10 observables in e^+e^-
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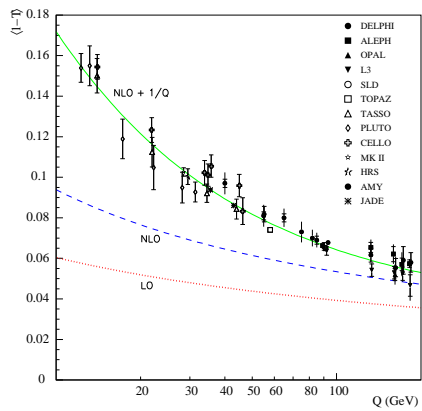
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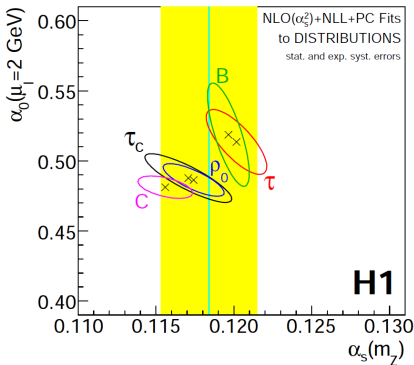
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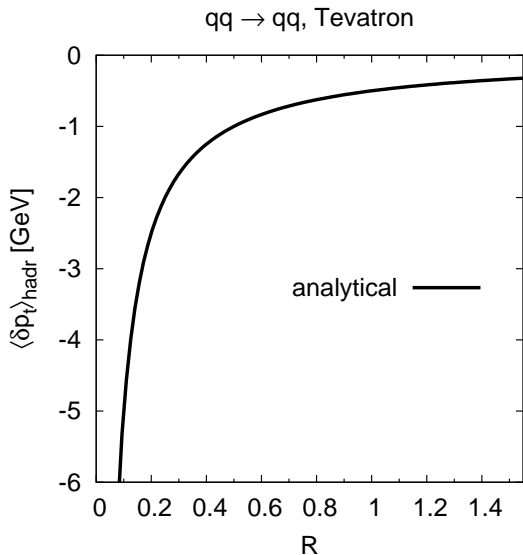
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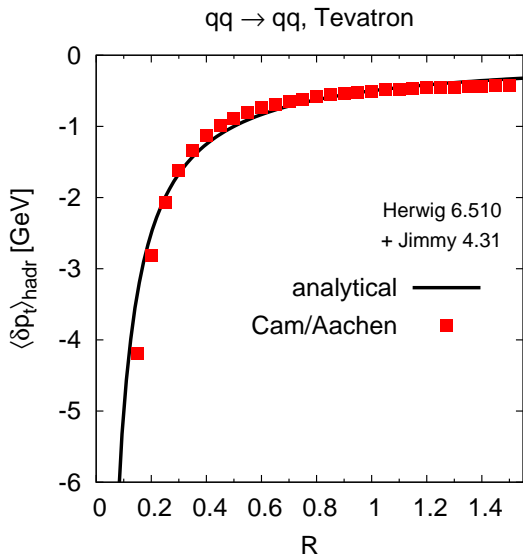
MC hadr. agrees with calc.

- ▶ to varying degrees for range of algs
- ▶ also in larger gluonic channels

MC UE \gg naive expectation

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- ▶ UE is huge at LHC
- ▶ largely indep. of scattering channel

Scale for (non-perturbative!) UE is ~ 10 GeV



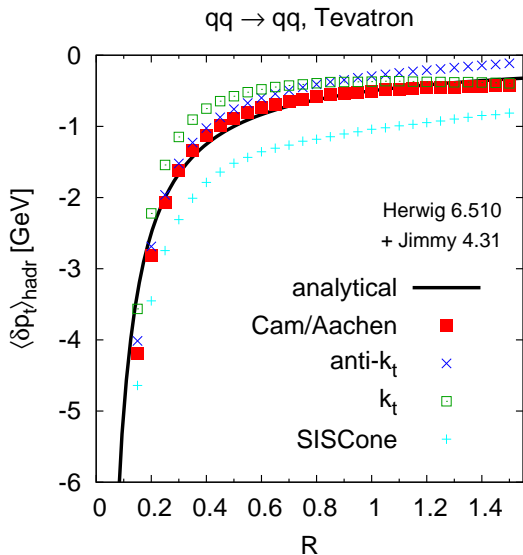
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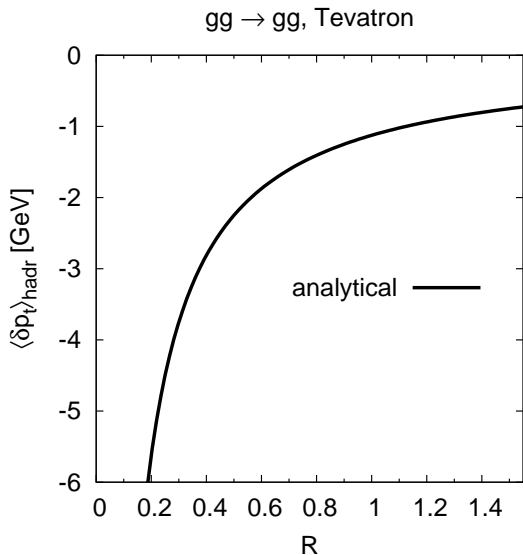
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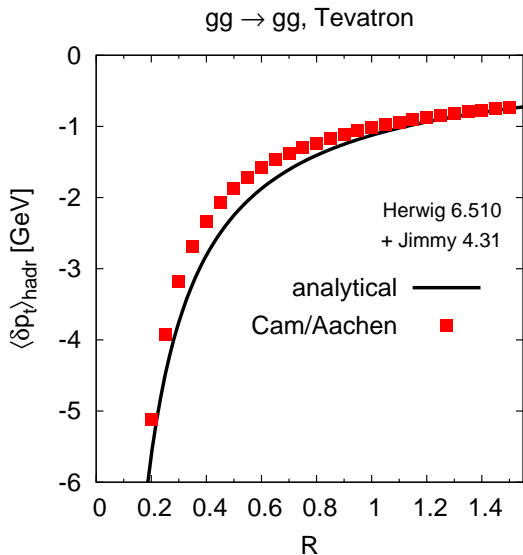
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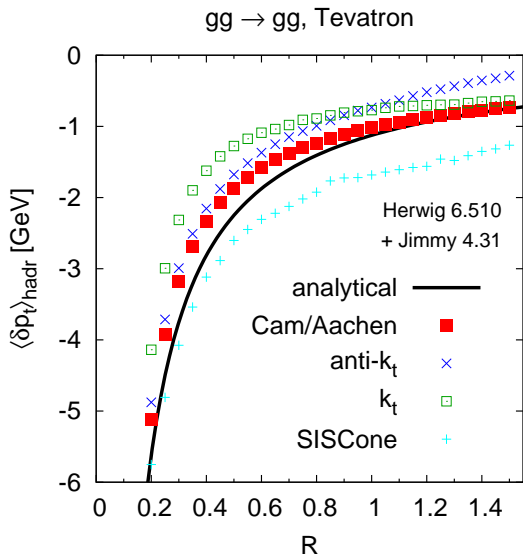
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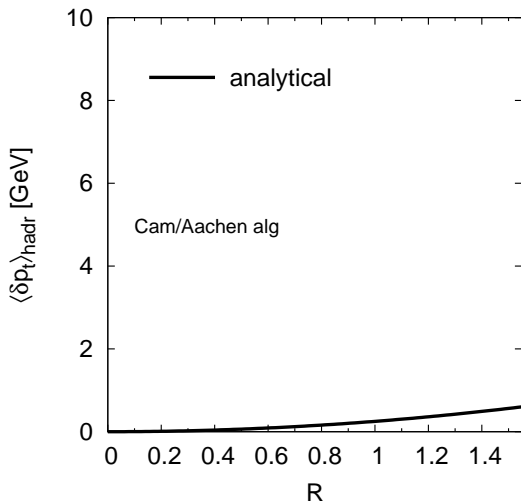
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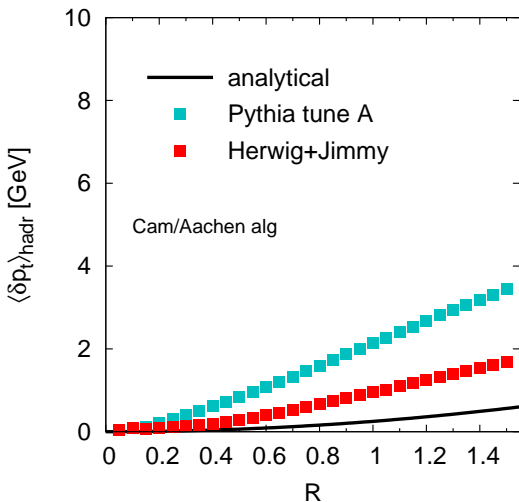
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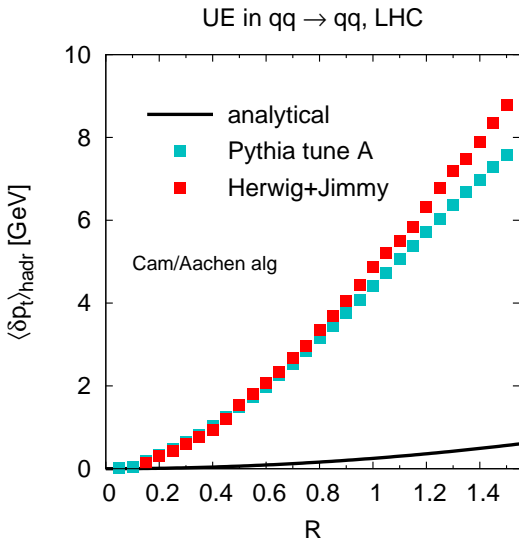
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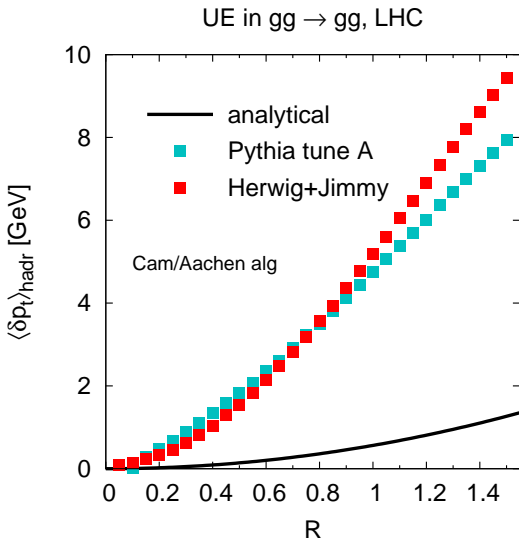
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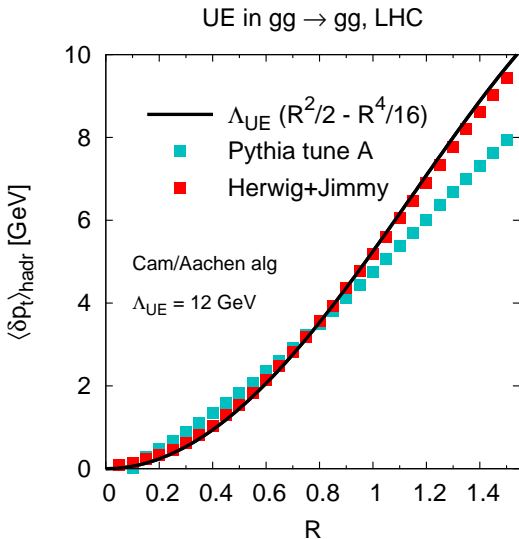
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To get best experimental resolutions, minimise contributions from all 3 components.

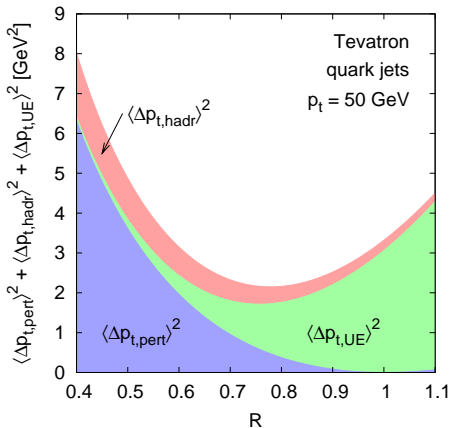
Here: sum of squared means
 Better still: calculate fluctuations

NB: this is rough picture
 details of p_t scaling wrong

But can still be used to understand general principles.

- 3. Understanding jet algs
- 2. Optimising parameters

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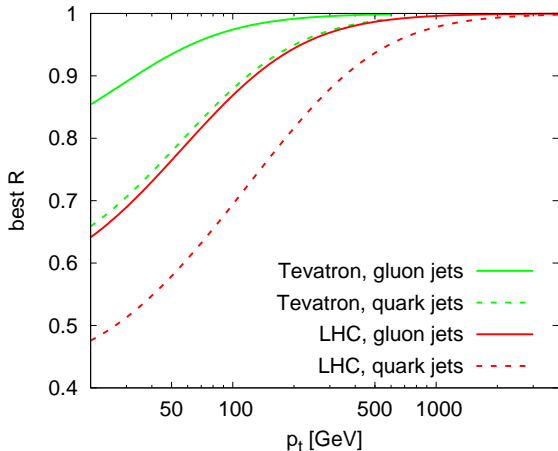
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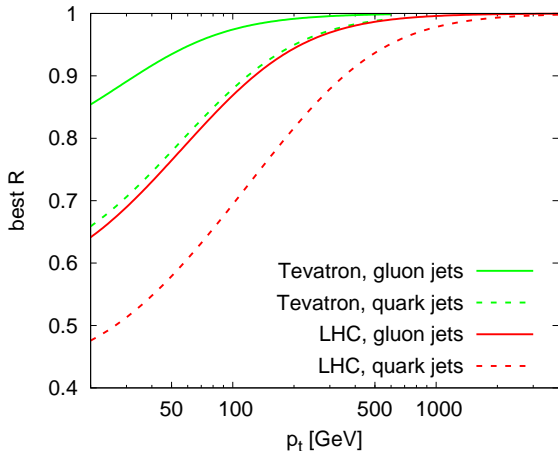
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Gluon jets wider
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This kind of information is the start of what might go into
“auto-focus for jetography”

This last part of talk was an overview of *1 of several* recent jet topics

Others include

- ▶ Subtraction of pileup Cacciari & GPS '07
- ▶ Jet areas \leftrightarrow sensitivity to UE/pileup Cacciari, GPS & Soyez prelim
- ▶ “Optimising R ” — cross checking with MC Cacciari, Rojo, GPS & Soyez, for Les Houches
- ▶ Jet flavour — e.g. reducing b -jet theory uncertainties from 40 – 60% to 10 – 20%. Banfi, GPS & Zanderighi '06, '07

- ▶ Jets are the closest we can get to seeing and giving meaning to partons
- ▶ Play a pivotal role in experimental analyses, comparisons to QCD calculations
- ▶ Significant progress in past 2 years towards making them *consistent* (IR/Collinear safe) and *practical* Link with computational geometry
All tools are made public:
<http://www.lpthe.jussieu.fr/~salam/fastjet/>
- ▶ The physics of how jets behave in a hadron-collider environment is a rich subject — much to be understood, and potential for significant impact in how jets are used at LHC